

Evaluating Soundness of a Gradual Verifier with Property Based Testing

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Software and Societal
Systems Department

Overview

Why Gradual Verification?

Static Verification

```
int findMax(Node l)
{
    int m = l→val;
    Node curr = l→next;
    while(curr ≠ null) {
        if(curr→val > m) {
            m = curr→val;
        }
        curr = curr→next;
    }
    return m;
}
```

Static Verification: Does not support incrementality

```
int findMax(Node l)
  requires l ≠ NULL
  ensures max(result,l) && contains(result,l)
{
  int m = l→val;
  Node curr = l→next;
  FOLDS/UNFOLDS
  while(curr ≠ null) { LOOP INVARIANTS
    if(curr→val > m) {
      m = curr→val;
      FOLDS/UNFOLDS
      LEMMAS
    }
    FOLDS/UNFOLDS
    curr = curr→next;
  }
  return m;
}
```

Dynamic Verification

```
int findMax(Node l)
  ensures max(result,l) && contains(result,l)
{
  int m = l→val;
  Node curr = l→next;
  while(curr ≠ null) {
    if(curr→val > m) {
      m = curr→val;
    }
    curr = curr→next;
  }
  return m;
}
```

Dynamic Verification: Runtime overhead is too much

```
int findMax(Node l)
```

```
  ens
```

```
{
```

```
  int
```

```
  Node
```

```
  while
```

```
  {
```

```
    i
```

```
  }
```

```
  C
```

```
}
```

```
ret
```

```
}
```

yarnpkg/berry

#1817 Benchmark PnP runtime overhead



43 comments



remorses opened on September 8, 2020



Gradual Verification
supports
incrementality.

Allows users to
specify as much as
they want.

Provides a **formal**
guarantee of
verifiability.

```
int findMax(Node l)
  requires ?
  ensures max(result,l) &&
  contains(result,l)
{
  int m = l→val;
  Node curr = l→next;
  while(curr ≠ NULL) ? {
    if(curr→val > m) {
      m = curr→val;
    }
    curr = curr→next;
  }

  return m;
}
```

Gradual Verification

```
int findMax(Node l)
  requires ? && l ≠ NULL
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{
  int m = l→val;
  Node curr = l→next;
  while(curr ≠ NULL) ? && LOOP INVARIANTS {
    if(curr→val > m) {
      m = curr→val;
    }
    curr = curr→next;
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```


Gradual Verification: Formal guarantee of verifiability

```
int findMax(Node l)
  requires ? && l ≠ NULL
  ensures max(result,l) && contains(result,l)
{
  int m = l→val;
  Node curr = l→next;
  FOLDS/UNFOLDS
  while(curr ≠ NULL) FOLDS/UNFOLDS {
    if(curr→val > m) {
      m = curr→val;
      FOLDS/UNFOLDS
      LEMMAS
    }
    FOLDS/UNFOLDS
    curr = curr→next;
  }

  return m;
}
```

Gradual Guarantee:
Verifiability and
reducibility are
**monotone with
respect to precision.**

*Conservative
extension:* Anything
provable in the
statically should be
provable in the
gradually.

Gradual C0: Design
has been proven
sound.

Extended with Gradual Formulas

**Static Verifier with
Implicit Dynamic
Frames and recursive
Abstract predicates**

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Lifting (Garcia et al. '16)

**Optimistic
Static
Verification
System**

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Lifting (Garcia et al. '16)

Optimistic
Static
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System

Dynamic
Verification
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Goal

Ensure Gradual CO's implementation is sound and allow for scalable bug fixes.

Why Property Based Testing?

A number of bugs had been **caught by hand**, in which *Gradual CO's* design was unsound



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Why Property Based Testing?

A number of bugs had been **caught by hand**, in which *Gradual CO's* design was **unsound**

There are no **lightweight** techniques available

Capturing the *truthiness* of a property's result provides **good coverage** for finding these implementation bugs.



Three-stage Pipeline

Reference model language
**Uses Gradual C0's
specification language**



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WIP: Input Generator
**Test suite of examples
that are not supposed to
verify correctly**



Three-stage Pipeline

Reference model language
Uses Gradual C0's
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WIP: Input Generator
Test suite of examples
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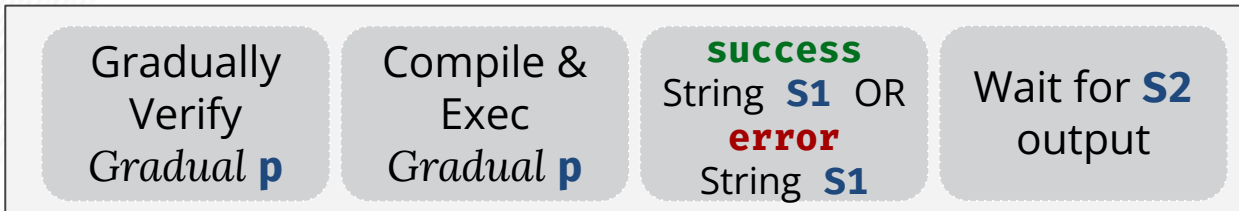
Checker: Dynamic C0
Asserts runtime checks
everywhere.
The ground truth

Checker: Gradual C0

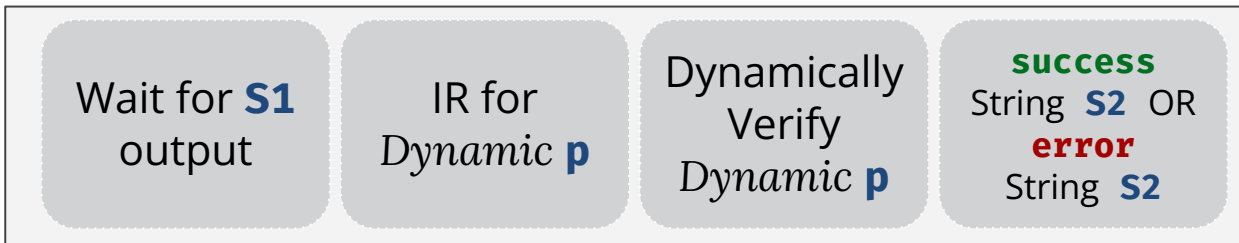


Checker Architecture

Input
C0 program **p**



Intermediate
Representation
for *Gradual* **p**



Output
if **S1** \neq **S2** and
S1 = **success**
then : **error**
else : **success**

By Example
**Evaluating Soundness of
Gradual C0 using PBT**

To prevent trivial failure,
we must **avoid specifying
preconditions and
folds/unfolds** that won't
be met while running

The truthiness for all
programs consists of a
pair of file executions

Binary Search Tree
breaking order:
left node > right node

```
- if (x < v) {  
+ if (v < x) {  
    if (l ≠ NULL) {  
        root→left = tree_add_helper(l,  
x, min, v-1);  
    } else {  
        root→left =  
create_tree_helper(x, min, v-1);  
    }  
    } else {  
- if (v < x) {  
+ if (x < v) {  
    if (r ≠ NULL) {  
        root→right =  
tree_add_helper(r, x, v+1, max);  
    } else {  
        root→right =  
create_tree_helper(x, v+1, max);  
    }  
    }  
}
```

Input Generator Architecture

We caught **4 soundness bugs** at different implementation phases of *Gradual C0*

```
/*@
predicate list(struct Node *l) =
  ? && (l ≠ NULL ? acc(l→value)
&& list(l→next) : true);
@*/
void append(Node *root, int value)
  //@requires ?;
  //@ensures ? && list(root);
{
  Node *n = root;
  while (n→next ≠ NULL)
    //@loop_invariant ?;
    n = n→next;
  n→next = alloc(Node);
  n→next→value = value;
}
```

Future Work

Carnegie Mellon University

School of Computer Science



S3D

Software and Societal
Systems Department

Check out our website to
learn more about our work:

S3D.cmu.edu

